

# Randomized Algorithms 2026 !

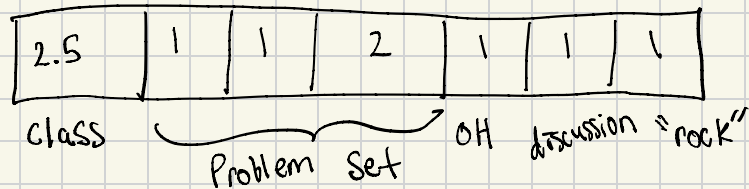
This class:

- probability + linear algebra
- data science at scale
- beautiful math

challenging course!

Prereqs: linear algebra, probability, algorithms

Time Expectation: 9.5 hr/week



- Earlier = better
- More drawing = better
- More talking = better

[www.rtealwitter.com/rads2026](http://www.rtealwitter.com/rads2026)

- discord for communication
- reading for each lecture
- these slides are online

Plan:

- Math Review
- Streaming and Sketching
- Linear Algebra and Spectral Methods

Course Goal: Use randomness to solve problems more efficiently

## Probability

Random variable  $X$

eg,  $X$  = outcome of fair dice

For  $x \in \{1, 2, 3, 4, 5, 6\}$ ,

$$\Pr(X=x) = 1/6$$

What is  $X$  on average?

Expectation of  $X$ :

$$\mathbb{E}[X] = \sum_x x \Pr(X=x)$$

$$= \sum_{x=1}^6 x \cdot 1/6$$

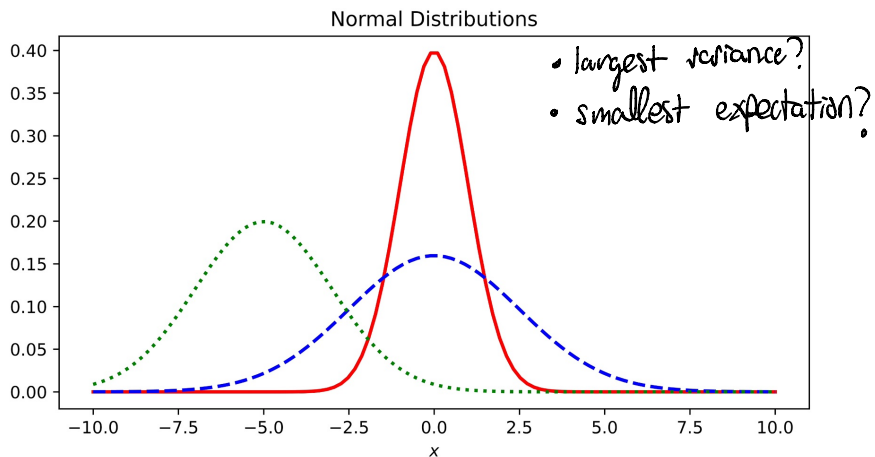
$$= 1/6 (1+2+\dots+6)$$

$$= 2\frac{1}{6}$$

How close is  $X$  to its expectation?

Variance of  $X$ :

$$\begin{aligned}\text{Var}(X) &= \mathbb{E}[(X - \mathbb{E}[X])^2] \\ &= \sum_x (x - \mathbb{E}[X])^2 \Pr(X=x) \\ &= \end{aligned}$$



## Multiplying by Scalars

$$\begin{aligned}\mathbb{E}[\alpha X] &= \sum_x \alpha x \Pr(X=x) \\ &= \alpha \sum_x x \Pr(X=x) \\ &= \alpha \mathbb{E}[X]\end{aligned}$$

$$\text{Var}(\alpha X) =$$

## Independent Random Variables

$A, B$  are events

$A$  = dice shows 1 or 2

$B$  = dice roll is odd

$$\begin{aligned} \overset{\text{"and"}}{Pr(A \cap B)} &= Pr(A) Pr(B | A) \\ &= Pr(B) Pr(A | B) \end{aligned}$$

$$Pr(B | A) = Pr(B) \quad \overset{\text{"if and only if"}}{\text{iff}} \quad A, B \text{ independent}$$

$$\Leftrightarrow Pr(A \cap B) = Pr(A) Pr(B) \quad \text{iff} \quad A, B \text{ independent}$$

Example  $A, B$  indep?

Random Variables  $X, Y$  are indep

$$\text{iff} \quad Pr(X=x \cap Y=y) = Pr(X=x) Pr(Y=y)$$

for all possible outcomes  $x, y$

Independent?

- Heads on 1<sup>st</sup> flip, Heads on 2<sup>nd</sup>
- Ace on 1<sup>st</sup> draw, Ace on 2<sup>nd</sup>
- Rain, carrying an umbrella



## Linearity of Expectation

Thm:  $E[X + Y] = E[X] + E[Y]$

Proof: 
$$\begin{aligned} E[X + Y] &= \sum_x \sum_y (x + y) \Pr(X=x \cap Y=y) \\ &= \sum_x \sum_y x \Pr(X=x \cap Y=y) + \sum_x \sum_y y \Pr(X=x \cap Y=y) \\ &= \sum_x x \sum_y \Pr(X=x \cap Y=y) + \sum_y y \sum_x \Pr(Y=y \cap X=x) \\ &= \sum_x x \Pr(X=x) + \sum_y y \Pr(Y=y) \\ &= E[X] + E[Y] \end{aligned}$$

Expected number of shared birthdays?

## Week 1 Thursday

- No OH Monday 1/26
- ⇒ Problem 2 due 2/2
- Join discord!
- Discussions!
- ↳ host can write quiz question

Today:

- More expectation / variance
- Set size estimation!

## Expectation & Variance

$$E[X+Y] = E[X] + E[Y]$$

- never?
- sometimes?
- always?

If  $X, Y$  indep,  $E[XY] = E[X]E[Y]$

$$E[XY] = \sum_{x,y} xy \Pr(X=x \cap Y=y)$$

indep  $\rightarrow$

$$= \sum_x \sum_y xy \Pr(X=x) \Pr(Y=y)$$

$$= \sum_x x \Pr(X=x) \sum_y y \Pr(Y=y)$$

$$= E[X] E[Y]$$

Dependent Example

$$X = \begin{cases} 1 & \text{if tails} \\ 0 & \text{if heads} \end{cases} \quad Y = \begin{cases} 1 & \text{if heads} \\ 0 & \text{if tails} \end{cases}$$

$$\text{Var}(X) = E[X^2] - E[X]^2$$

$$\text{Var}(X) = E[(X - E[X])^2]$$

$$= E[X^2 - 2XE[X] + E[X]^2]$$

$$= E[X^2] - 2E[X]E[X] + E[X]^2$$

$$= E[X^2] - E[X]^2$$

If  $X, Y$  indep,  $\text{Var}(X+Y) = \text{Var}(X) + \text{Var}(Y)$

← New person writes  
please!

## Set Size Estimation

Problem:

Given uniform query access to a set, estimate the size of the set.

- Ecology
- social networks
- Internet indexing

Suppose Wikipedia claims it has

$n = 1,000,000$  articles, how do we check?

Naive Solution: keep querying until we get  $n$  unique elements

The issue is that the expected queries to find all  $n$  is  $n \log n$

Clever Solution: Count duplicates!

Intuitively, more duplicates  $\rightarrow$  smaller  $n$

## Coupon Collector Problem

$T = \#$  queries to get all  $n$

$$= T_1 + T_2 + \dots + T_n$$

← queries to get  $n$ th unique element

← waiting to get success  
 $T_i \sim \text{Geometric}(p_i)$

with  $p_i = \frac{n - (i-1)}{n}$

$$\mathbb{E}[T_i] = \sum_{k=1}^{\infty} (1-p_i)^{k-1} p_i = p_i \frac{1}{(1-(1-p_i))^2} = \frac{1}{p_i} = \frac{n}{n-(i-1)}$$

$$\mathbb{E}[T] = \sum_{i=1}^n \mathbb{E}[T_i] = \sum_{i=1}^n \frac{n}{n-(i-1)} = n \left( \frac{1}{n} + \frac{1}{n-1} + \dots + \frac{1}{1} \right) = n \sum_{j=1}^n \frac{1}{j} \approx n \log n$$

Harmonic series

$$\frac{\partial}{\partial x} \sum_{k=0}^{\infty} x^k = \frac{\partial}{\partial x} \frac{1}{1-x}$$
$$\sum_{k=1}^{\infty} k x^{k-1} = \frac{1}{(1-x)^2}$$

Define

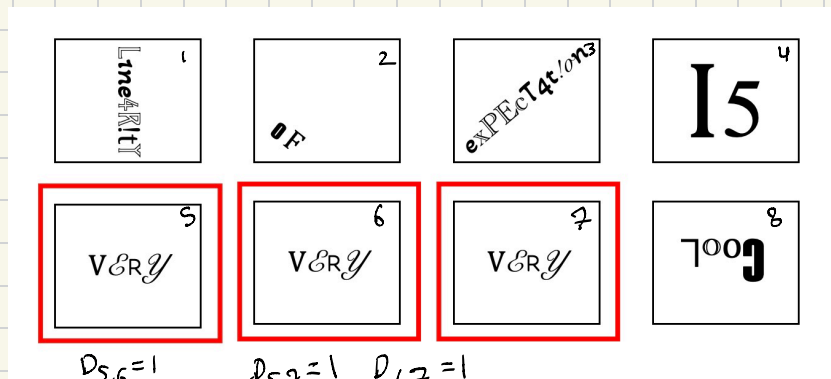
$$D_{i,j} = \begin{cases} 1 & \text{if } i \text{ and } j \text{ same} \\ 0 & \text{else} \end{cases}$$

$$D = \sum_{i,j=1 \leq i < j \leq m} D_{i,j}$$

$$\mathbb{E}[D] = \sum_{i < j} \mathbb{E}[D_{i,j}] = \binom{m}{2} \frac{1}{n} = \frac{m(m-1)}{2n}$$

Suppose we made  $m=1000$  queries and saw  $D=10$  duplicates.

How does this compare to what we expect?



## Markov's Inequality

Theorem: For any non-negative rv  $X$  and  $t > 0$ ,

$$\Pr(X \geq t) \leq \frac{E[X]}{t}$$

Proof:  $E[X] = \sum_x x \Pr(X=x)$

$$= \sum_{x: x \geq t} x \Pr(X=x) + \sum_{x: x < t} x \Pr(X=x)$$
$$\geq \sum_{x: x \geq t} t \Pr(X=x) + 0 = t \Pr(X \geq t)$$

Answer to duplicate question:

$$\Pr(D \geq 10) \leq \frac{E[D]}{10} = \frac{.4995}{10} = .04995$$